Chapter 20

Low-Level Programming



Copyright © 2008 W. W. Norton & Company. All rights reserved.

1

Introduction

- Previous chapters have described C's high-level, machine-independent features.
- However, some kinds of programs need to perform operations at the bit level:
 - Systems programs (including compilers and operating systems)
 - Encryption programs
 - Graphics programs
 - Programs for which fast execution and/or efficient use of space is critical



Bitwise Operators

- C provides six *bitwise operators*, which operate on integer data at the bit level.
- Two of these operators perform shift operations.
- The other four perform bitwise complement, bitwise *and*, bitwise exclusive *or*, and bitwise inclusive *or* operations.



Bitwise Shift Operators

- The bitwise shift operators shift the bits in an integer to the left or right:
 - << left shift
 - >> right shift
- The operands for << and >> may be of any integer type (including char).
- The integer promotions are performed on both operands; the result has the type of the left operand after promotion.



Bitwise Shift Operators

- The value of $i \ll j$ is the result when the bits in
 - i are shifted left by j places.
 - For each bit that is "shifted off" the left end of i, a zero bit enters at the right.
- The value of i >> j is the result when i is shifted right by j places.
 - If i is of an unsigned type or if the value of i is nonnegative, zeros are added at the left as needed.
 - If i is negative, the result is implementation-defined.



Bitwise Shift Operators

• Examples illustrating the effect of applying the shift operators to the number 13: unsigned short i, j;



Bitwise Shift Operators

- To modify a variable by shifting its bits, use the compound assignment operators <<= and >>=:
 - i = 13;
 /* i is now 13 (binary 000000000001101) */
 - i <<= 2; /* i is now 52 (binary 000000000110100) */
 - i >>= 2;
 /* i is now 13 (binary 000000000001101) */



Bitwise Shift Operators

- The bitwise shift operators have lower precedence than the arithmetic operators, which can cause surprises:
 - i << 2 + 1 means i << (2 + 1), not (i << 2) + 1



Bitwise Complement, *And*, Exclusive *Or*, and Inclusive *Or*

- There are four additional bitwise operators:
 - ~ bitwise complement
 - & bitwise *and*
 - ^ bitwise exclusive or
 - | bitwise inclusive or
- The ~ operator is unary; the integer promotions are performed on its operand.
- The other operators are binary; the usual arithmetic conversions are performed on their operands.



Bitwise Complement, *And*, Exclusive *Or*, and Inclusive *Or*

- The ~, &, ^, and | operators perform Boolean operations on all bits in their operands.
- The ^ operator produces 0 whenever both operands have a 1 bit, whereas | produces 1.



Bitwise Complement, *And*, Exclusive *Or*, and Inclusive *Or*

• Examples of the ~, &, ^, and | operators: unsigned short i, j, k; i = 21;/* i is now 21 (binary 0000000000010101) */ i = 56;/* j is now 56 (binary 000000000111000) */ $k = \sim i$: /* k is now 65514 (binary 1111111111101010) */ k = i & j;/* k is now 16 (binary 000000000000000) */ $k = i ^{\prime} j;$ /* k is now 45 (binary 000000000101101) */ k = i | j;61 (binary 000000000111101) */ /* k is now PROGRAMMING Copyright © 2008 W. W. Norton & Company. 11 All rights reserved. A Modern Approach Second Edition

Bitwise Complement, *And*, Exclusive *Or*, and Inclusive *Or*

- The ~ operator can be used to help make low-level programs more portable.
 - An integer whose bits are all $1: \sim 0$
 - An integer whose bits are all 1 except for the last five:
 ~0x1f



Bitwise Complement, *And*, Exclusive *Or*, and Inclusive *Or*

Each of the ~, &, ^, and | operators has a different precedence:

Highest: ~

&

Lowest:

- Examples:
 - i&~j|k means (i&(~j))|k
 - i^j&~k means i^ (j& (~k))
- Using parentheses helps avoid confusion.



Bitwise Complement, *And*, Exclusive *Or*, and Inclusive *Or*

The compound assignment operators &=, ^=, and |
 = correspond to the bitwise operators &, ^, and |:



- The bitwise operators can be used to extract or modify data stored in a small number of bits.
- Common single-bit operations:
 - Setting a bit
 - Clearing a bit
 - Testing a bit
- Assumptions:
 - i is a 16-bit unsigned short variable.
 - The leftmost—or *most significant*—bit is numbered 15 and the least significant is numbered 0.



• *Setting a bit.* The easiest way to set bit 4 of *i* is to *or* the value of *i* with the constant 0x0010:

• If the position of the bit is stored in the variable j, a shift operator can be used to create the mask:

i |= 1 << j; /* sets bit j */

 Example: If j has the value 3, then 1 << j is 0x0008.



- *Clearing a bit.* Clearing bit 4 of *i* requires a mask with a 0 bit in position 4 and 1 bits everywhere else:
 - i = 0x00ff;
 /* i is now 000000011111111 */
 - i &= ~0x0010; /* i is now 000000011101111 */
- A statement that clears a bit whose position is stored in a variable:



• *Testing a bit.* An if statement that tests whether bit 4 of i is set:

if (i & 0x0010) ... /* tests bit 4 */

A statement that tests whether bit j is set:
 if (i & 1 << j) ... /* tests bit j */



- Working with bits is easier if they are given names.
- Suppose that bits 0, 1, and 2 of a number correspond to the colors blue, green, and red, respectively.
- Names that represent the three bit positions:

#define BLUE 1
#define GREEN 2
#define RED 4



- Examples of setting, clearing, and testing the BLUE bit:
 - i |= BLUE; /* sets BLUE bit */
 i &= ~BLUE; /* clears BLUE bit */
 i (* bit */
 - if (i & BLUE) ... /* tests BLUE bit */



• It's also easy to set, clear, or test several bits at time:

/* tests BLUE and GREEN bits */

• The if statement tests whether either the BLUE bit or the GREEN bit is set.



- Dealing with a group of several consecutive bits (a *bit-field*) is slightly more complicated than working with single bits.
- Common bit-field operations:
 - Modifying a bit-field
 - Retrieving a bit-field



- *Modifying a bit-field*. Modifying a bit-field requires two operations:
 - A bitwise *and* (to clear the bit-field)
 - A bitwise *or* (to store new bits in the bit-field)
- Example:
 - $i = i \& \sim 0 \times 0070 | 0 \times 0050;$
 - /* stores 101 in bits 4-6 */
- The & operator clears bits 4–6 of i; the | operator then sets bits 6 and 4.



- To generalize the example, assume that j contains the value to be stored in bits 4–6 of i.
- j will need to be shifted into position before the bitwise *or* is performed:

• The | operator has lower precedence than & and <<, so the parentheses can be dropped:

i = i & ~0x0070 | j << 4;



• *Retrieving a bit-field.* Fetching a bit-field at the right end of a number (in the least significant bits) is easy:

• If the bit-field isn't at the right end of i, we can first shift the bit-field to the end before extracting the field using the & operator:



Program: XOR Encryption

- One of the simplest ways to encrypt data is to exclusive-*or* (XOR) each character with a secret key.
- Suppose that the key is the & character.
- XORing this key with the character z yields the \ character:

 00100110
 (ASCII code for &)

 XOR
 01111010
 (ASCII code for z)

 01011100
 (ASCII code for \)



Program: XOR Encryption

• Decrypting a message is done by applying the same algorithm:



Program: XOR Encryption

- The xor.c program encrypts a message by XORing each character with the & character.
- The original message can be entered by the user or read from a file using input redirection.
- The encrypted message can be viewed on the screen or saved in a file using output redirection.



Program: XOR Encryption

• A sample file named msg:

Trust not him with your secrets, who, when left alone in your room, turns over your papers. --Johann Kaspar Lavater (1741-1801)

• A command that encrypts msg, saving the encrypted message in newmsg:

xor <msg >newmsg

• Contents of newmsg:

rTSUR HIR NOK QORN _IST UCETCRU, QNI, QNCH JC@R GJIHC OH _IST TIIK, RSTHU IPCT _IST VGVCTU. --linghh mGUVGT jGPGRCT (1741-1801)



Program: XOR Encryption

• A command that recovers the original message and displays it on the screen:

xor <newmsg</pre>



Program: XOR Encryption

- The xor.c program won't change some characters, including digits.
- XORing these characters with & would produce invisible control characters, which could cause problems with some operating systems.
- The program checks whether both the original character and the new (encrypted) character are printing characters.
- If not, the program will write the original character instead of the new character.



xor.c

```
/* Performs XOR encryption */
#include <ctype.h>
#include <stdio.h>
#define KEY '&'
int main (void)
{
  int orig char, new char;
 while ((orig char = getchar()) != EOF) {
    new char = orig char ^ KEY;
    if (isprint(orig char) && isprint(new char))
      putchar(new char);
    else
      putchar(orig char);
  }
```

return 0;



Bit-Fields in Structures

- The bit-field techniques discussed previously can be tricky to use and potentially confusing.
- Fortunately, C provides an alternative: declaring structures whose members represent bit-fields.



Bit-Fields in Structures

- Example: How DOS stores the date at which a file was created or last modified.
- Since days, months, and years are small numbers, storing them as normal integers would waste space.
- Instead, DOS allocates only 16 bits for a date, with 5 bits for the day, 4 bits for the month, and 7 bits for the year:





34

Bit-Fields in Structures

• A C structure that uses bit-fields to create an identical layout:

```
struct file_date {
    unsigned int day: 5;
    unsigned int month: 4;
    unsigned int year: 7;
};
```

• A condensed version:

```
struct file_date {
    unsigned int day: 5, month: 4, year: 7;
};
```



Bit-Fields in Structures

- The type of a bit-field must be either int, unsigned int, or signed int.
- Using int is ambiguous; some compilers treat the field's high-order bit as a sign bit, but others don't.
- In C99, bit-fields may also have type _Bool.
- C99 compilers may allow additional bit-field types.



Bit-Fields in Structures

• A bit-field can be used in the same way as any other member of a structure:

```
struct file_date fd;
```

• Appearance of the fd variable after these assignments:





37

Bit-Fields in Structures

- The address operator (&) can't be applied to a bit-field.
- Because of this rule, functions such as scanf can't store data directly in a bit-field: scanf("%d", &fd.day); /*** WRONG ***/
- We can still use scanf to read input into an ordinary variable and then assign it to fd.day.



- The C standard allows the compiler considerable latitude in choosing how it stores bit-fields.
- The rules for handling bit-fields depend on the notion of "storage units."
- The size of a storage unit is implementation-defined.
 - Typical values are 8 bits, 16 bits, and 32 bits.



- The compiler packs bit-fields one by one into a storage unit, with no gaps between the fields, until there's not enough room for the next field.
- At that point, some compilers skip to the beginning of the next storage unit, while others split the bit-field across the storage units.
- The order in which bit-fields are allocated (left to right or right to left) is also implementation-defined.



- Assumptions in the file_date example:
 - Storage units are 16 bits long.
 - Bit-fields are allocated from right to left (the first bitfield occupies the low-order bits).
- An 8-bit storage unit is also acceptable if the compiler splits the month field across two storage units.



- The name of a bit-field can be omitted.
- Unnamed bit-fields are useful as "padding" to ensure that other bit-fields are properly positioned.
- A structure that stores the time associated with a DOS file:

```
struct file_time {
    unsigned int seconds: 5;
    unsigned int minutes: 6;
    unsigned int hours: 5;
};
```



How Bit-Fields Are Stored

• The same structure with the name of the seconds field omitted:

• The remaining bit-fields will be aligned as if seconds were still present.



How Bit-Fields Are Stored

• The length of an unnamed bit-field can be 0:

```
struct s {
    unsigned int a: 4;
    unsigned int : 0;    /* 0-length bit-field */
    unsigned int b: 8;
};
```

- A 0-length bit-field tells the compiler to align the following bit-field at the beginning of a storage unit.
 - If storage units are 8 bits long, the compiler will allocate 4 bits for a, skip 4 bits to the next storage unit, and then allocate 8 bits for b.
 - If storage units are 16 bits long, the compiler will allocate 4 bits for a, skip 12 bits, and then allocate 8 bits for b.



Other Low-Level Techniques

- Some features covered in previous chapters are used often in low-level programming.
- Examples:
 - Defining types that represent units of storage
 - Using unions to bypass normal type-checking
 - Using pointers as addresses
- The volatile type qualifier was mentioned in Chapter 18 but not discussed because of its lowlevel nature.



Defining Machine-Dependent Types

- The char type occupies one byte, so characters can be treated as bytes.
- It's a good idea to define a BYTE type: typedef unsigned char BYTE;
- Depending on the machine, additional types may be needed.
- A useful type for the x86 platform: typedef unsigned short WORD;



- Unions can be used in a portable way, as shown in Chapter 16.
- However, they're often used in C for an entirely different purpose: viewing a block of memory in two or more different ways.
- Consider the file_date structure described earlier.
- A file_date structure fits into two bytes, so any two-byte value can be thought of as a file_date structure.



- In particular, an unsigned short value can be viewed as a file_date structure.
- A union that can be used to convert a short integer to a file date or vice versa:

```
union int_date {
    unsigned short i;
    struct file_date fd;
};
```



• A function that prints an unsigned short argument as a file date:

```
void print_date(unsigned short n)
{
    union int_date u;
```



- Using unions to allow multiple views of data is especially useful when working with registers, which are often divided into smaller units.
- x86 processors have 16-bit registers named AX, BX, CX, and DX.
- Each register can be treated as two 8-bit registers.
 - AX is divided into registers named AH and AL.



- Writing low-level applications for x86-based computers may require variables that represent AX, BX, CX, and DX.
- The goal is to access both the 16- and 8-bit registers, taking their relationships into account.
 - A change to AX affects both AH and AL; changing AH or AL modifies AX.
- The solution is to set up two structures:
 - The members of one correspond to the 16-bit registers.
 - The members of the other match the 8-bit registers.



• A union that encloses the two structures:

```
union {
  struct {
    WORD ax, bx, cx, dx;
  } word;
  struct {
    BYTE al, ah, bl, bh, cl, ch, dl, dh;
  } byte;
}
```

```
} regs;
```



- The members of the word structure will be overlaid with the members of the byte structure.
 ax will occupy the same memory as al and ah.
- An example showing how the regs union might be used:

```
regs.byte.ah = 0x12;
regs.byte.al = 0x34;
printf("AX: %hx\n", regs.word.ax);
```

• Output:

AX: 1234



- Note that the byte structure lists al before ah.
- When a data item consists of more than one byte, there are two logical ways to store it in memory:
 - *Big-endian:* Bytes are stored in "natural" order (the leftmost byte comes first).
 - *Little-endian:* Bytes are stored in reverse order (the leftmost byte comes last).
- x86 processors use little-endian order.



- We don't normally need to worry about byte ordering.
- However, programs that deal with memory at a low level must be aware of the order in which bytes are stored.
- It's also relevant when working with files that contain non-character data.



Using Pointers as Addresses

- An address often has the same number of bits as an integer (or long integer).
- Creating a pointer that represents a specific address is done by casting an integer to a pointer:
 BYTE *p;



Program: Viewing Memory Locations

- The viewmemory.c program allows the user to view segments of computer memory.
- The program first displays the address of its own main function as well as the address of one of its variables.
- The program next prompts the user to enter an address (as a hexadecimal integer) plus the number of bytes to view.
- The program then displays a block of bytes of the chosen length, starting at the specified address.



Program: Viewing Memory Locations

- Bytes are displayed in groups of 10 (except for the last group).
- Bytes are shown both as hexadecimal numbers and as characters.
- Only printing characters are displayed; other characters are shown as periods.
- The program assumes that int values and addresses are stored using 32 bits.
- Addresses are displayed in hexadecimal.



viewmemory.c

/* Allows the user to view regions of computer memory */

```
#include <ctype.h>
#include <stdio.h>
```

```
typedef unsigned char BYTE;
```

```
int main(void)
{
    unsigned int addr;
    int i, n;
```

```
BYTE *ptr;
```

printf("Address of main function: %x\n", (unsigned int) main);
printf("Address of addr variable: %x\n", (unsigned int)
 &addr);





```
ptr = (BYTE *) addr;
for (; n > 0; n -= 10) {
  printf("%8X ", (unsigned int) ptr);
  for (i = 0; i < 10 \&\& i < n; i++)
   printf("%.2X ", *(ptr + i));
  for (; i < 10; i++)
   printf(" ");
  printf(" ");
  for (i = 0; i < 10 \&\& i < n; i++) {
    BYTE ch = *(ptr + i);
    if (!isprint(ch))
      ch = '.';
    printf("%c", ch);
  printf("\n");
 ptr += 10;
}
```

return 0;

}



Program: Viewing Memory Locations

• Sample output using GCC on an x86 system running Linux: Address of main function: 804847c

Address of addr variable: bff41154

Enter a (hex) address: <u>8048000</u> Enter number of bytes to view: <u>40</u>

Address					Byt	ces					Characters
8048000	<u></u> 7f	45	4C	46	01	01	01	00	00	00	.ELF
804800A	00	00	00	00	00	00	02	00	03	00	
8048014	01	00	00	00	С0	83	04	08	34	00	
804801E	00	00	С0	0A	00	00	00	00	00	00	• • • • • • • • • •

• The 7F byte followed by the letters E, L, and F identify the format (ELF) in which the executable file was stored.



Program: Viewing Memory Locations

• A sample that displays bytes starting at the address of addr: Address of main function: 804847c

Address of addr variable: bfec5484

Enter a (hex) address: <u>bfec5484</u> Enter number of bytes to view: <u>64</u>

Address					Byt	ces					Characters
BFEC5484	84	54	EC	BF	в0	54	EC	BF	F4	6F	.TTo
BFEC548E	68	00	34	55	ЕC	BF	С0	54	ЕC	BF	h.4UT
BFEC5498	08	55	ЕC	BF	ЕЗ	ЗD	57	00	00	00	.U=W
BFEC54A2	00	00	AO	BC	55	00	8 0	55	ЕC	BF	UU
BFEC54AC	ЕЗ	3D	57	00	01	00	00	00	34	55	.=W4U
BFEC54B6	ЕC	ΒF	ЗC	55	ЕC	BF	56	11	55	00	<uv.u.< td=""></uv.u.<>
BFEC54C0	F4	6F	68	00							.oh.

• When reversed, the first four bytes form the number BFEC5484, the address entered by the user.



- On some computers, certain memory locations are "volatile."
- The value stored at such a location can change as a program is running, even though the program itself isn't storing new values there.
- For example, some memory locations might hold data coming directly from input devices.



The volatile Type Qualifier

- The volatile type qualifier allows us to inform the compiler if any of the data used in a program is volatile.
- volatile typically appears in the declaration of a pointer variable that will point to a volatile memory location:

volatile BYTE *p;
 /* p will point to a volatile byte */



- Suppose that p points to a memory location that contains the most recent character typed at the user's keyboard.
- A loop that obtains characters from the keyboard and stores them in a buffer array:

```
while (buffer not full) {
    wait for input;
    buffer[i] = *p;
    if (buffer[i++] == '\n')
        break;
}
```



- A sophisticated compiler might notice that this loop changes neither p nor *p.
- It could optimize the program by altering it so that *p is fetched just once:

```
store *p in a register;
while (buffer not full) {
    wait for input;
    buffer[i] = value stored in register;
    if (buffer[i++] == '\n')
        break;
}
```



- The optimized program will fill the buffer with many copies of the same character.
- Declaring that p points to volatile data avoids this problem by telling the compiler that *p must be fetched from memory each time it's needed.

